Speeding up computation with a filesystem of network RAM Colin Schimmelfing '10, Advisor: Tia Newhall ~ Swarthmore College, Swarthmore PA

What is Nswap?



- When a computer runs out of RAM, it tries to 'swap' to
- 'Swapping' treats the disk as if it was RAM
- Problem: hard disks are one million times slower than F
- Storing data across the network in idle RAM is faster that
- Nswap allows a computer to 'borrow' RAM from other computers on the network
- Transparent to the program, handled at operating system (Programmer does not need to do anything to use N

Implementation

• Created a temporary filesystem that resides on the idle RAN computers in the local network

Needs to support multiple modes:

- **Swapping mode**, with page-size requests and garbage co • GC is needed to free up memory for future use from pro have finished- RAM can be used again
- Filesystem mode, with smaller requests and no GC
 - Can't use GC- no applications claim files, so GC thinks t be collected

• One problem: we don't want to use too much memory to ke where our pages are.



	Adding an API
	• Application Programming Interface (AP interaction with Nswap
disk RAM an disk em level Jswap)	 •Why do we want this? •More control to programmers can increase characteristics of their own data •Specify which elements of memory could be e. Could provide greater resources to program. • Courrently programmers have no control; New Currently pro
	Experime
M of	Comparing filesystem performance <u>Experiment 1: Direct write to files</u> • Simple writes and reads to files on disk an
ollection (GC) ograms that	• Should be best case for disk compared to I
that files can	 <u>Experiment 2: External merge sort perf</u> External merge sort allows sorting withou of the data in memory at a time
keep track of	• As long as two of the smallest units can fit memory, the sort can be completed
-size structures	 Requires many reads and writes to disk
etures have KB) in only as er chunks as we need vidth	 <u>Experiment 3: TPIE, a library for large 6</u> TPIE implements fast algorithms for data large to fit in memory Optimized for disk, sequential reads and w Ran the TPIE external merge sort algorith

to Nswap

- 'I) allows programmer direct
- performance- they know the
- e farther away (on the network) ns which try not to swap Iswap is transparent

- ust be written and read repeatedly
- faster access times
- cause some programs will limit се
- computer can not distinguish

ents

e between Nswap and disk

- nd files on our Nswap filesystem
- Nswap
- <u>tormance</u>
- it all
- in
- data sets sets too
- writes
- nm

<u>Merge Sort</u> 4 3 2 1 2 1 3 1 2 3

Experiments

- Discover a fix to TPIE issue
- memory datasets

<u>New Interfaces</u>

- Malloc-like interface to Nswap

Other Improvements

- performance
- Add flash memory capability
- Work on persistence for the filesystem

Results

Simple Reads and Writes

• 390.625 MB of data (100,000 pages at 4 KB each)

- Simple, sequential reads and writes to a file
- Nswap is about ¹/₃ faster than disk

- 300 MB of data with 30 MB of memory allowed for the merge sort
- The data is stored either on disk or on Nswap-The memory is not swapped- only local memory
- TPIE is about ¹/₄ faster than Nswap!
- This seems impossible, since disk is slower even
- for large continuous writes, disk's best case

Future Work

• Try running with STXXL, another library for larger-than-

• Other memory interfaces, such as mmap • Provide an estimate of the network RAM available

• Tune Nswap parameters to maximize filesystem